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Electronic Companion—“Stochastic Sequential Decision-Making
with a Random Number of Jobs” by Alexander G. Nikolaev and
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Proofs of Theorems

EC.1. Proof of Theorem 1.

THEOREM 1. *Let π'^* be an optimal policy for AP. Then, an optimal policy for BP, π^* , is obtained using rules: (1) whenever a job arrives and π'^* assigns a worker with success rate zero, discard the job, (2) whenever a job arrives and π'^* assigns a worker with success rate $p > 0$, assign a worker with the same success rate.*

Proof: Note that π^* and π'^* have the same expected rewards for BP and AP, respectively. Suppose that π^* is not optimal for BP, and that there exists a policy π_{better} that yields a higher expected reward. Then define policy π'_{better} for AP as follows:

- whenever a job arrives and π_{better} discards it, assign a worker with success rate zero
- whenever a job arrives and π_{better} assigns a worker with success rate p , assign a worker with the same success rate.

Consider the reward earned in BP by making assignments for sequence $s = \{x_j\}_{j=1}^N$ according to policy π_{better} . Then, the same reward is earned in AP by making assignments for sequence $s' = \{\{x'_j\}_{j=1}^{N_{max}} : x'_j = x_j \text{ for } j = 1, 2, \dots, N \text{ and } x'_j = 0 \text{ for } j = N + 1, N + 2, \dots, N_{max}\}$ according to policy π'_{better} . Also, by (1),

$$P(\{X'_n > 0\} \cap \{X'_{n+1} = 0\}) = 1 \cdot \frac{\sum_{i=2}^{N_{max}} P_i}{\sum_{i=1}^{N_{max}} P_i} \cdot \frac{\sum_{i=3}^{N_{max}} P_i}{\sum_{i=2}^{N_{max}} P_i} \cdot \dots \cdot \frac{\sum_{i=n}^{N_{max}} P_i}{\sum_{i=n-1}^{N_{max}} P_i} \cdot \frac{P_n}{\sum_{i=n}^{N_{max}} P_i} = P_n,$$

which means that the probability of encountering sequence s' in AP is equal to the probability of encountering sequence s in BP. Therefore, π_{better} and π'_{better} yield the same expected rewards for BP and AP, respectively, and hence, π'_{better} is a better policy than π'^* , which contradicts the assumption that π'^* is optimal. □

EC.2. Proof of Theorem 3.

THEOREM 3. *Whenever job $n = 1, 2, \dots, N_{max} - 1$ arrives in AP, the optimal assignment policy is to assign the n^{th} job to the worker with the m^{th} highest success rate (available at the time the assignment decision has to be made) if x'_n lies in the m^{th} highest of the intervals, defined by the*

fixed breakpoints $E[Z_{1,n+1}^{N_{max}}|X'_n > 0]$, $E[Z_{2,n+1}^{N_{max}}|X'_n > 0]$, ..., $E[Z_{N_{max}-n,n+1}^{N_{max}}|X'_n > 0]$. These breakpoints are computed recursively,

$$\begin{aligned}
E[Z_{m,n+1}^{N_{max}}|X'_n > 0] &= \frac{\sum_{i=n+1}^{N_{max}} P(N=i)}{\sum_{i=n}^{N_{max}} P(N=i)} (F_{n+1}(E[Z_{m,n+2}^{N_{max}}|X'_{n+1} > 0])E[Z_{m,n+2}^{N_{max}}|X'_{n+1} > 0]) \\
&+ \int_{E[Z_{m,n+2}^{N_{max}}|X'_{n+1} > 0]}^{E[Z_{m-1,n+1}^{N_{max}}|X'_{n+1} > 0]} x dF_{n+1}(x) + (1 - F_{n+1}(E[Z_{m-1,n+1}^{N_{max}}|X'_{n+1} > 0]))E[Z_{m-1,n+1}^{N_{max}}|X'_{n+1} > 0].
\end{aligned} \tag{EC.1}$$

Proof: Theorem 2 can be used to establish an optimal policy for AP. Observe that for any $1 \leq m \leq N_{max} - n + 1$ and $1 \leq n \leq N_{max} - 1$, computing the expectation $E[Z_{m,n+1}^{N_{max}}|\mathcal{F}_n]$ (see Theorem 2) involves only the random variables X'_j , $j = n + 1, n + 2, \dots, N_{max}$. The distributions of these variables depend only on the value X'_n (which is known at the time $E[Z_{m,n+1}^{N_{max}}|\mathcal{F}_n]$ is computed), and not on $X'_1, X'_2, \dots, X'_{n-1}$. Therefore, the stochastic process $\{X'_j\}_{j=1}^{N_{max}}$ has the Markovian property, and hence, $E[Z_{m,n+1}^{N_{max}}|\mathcal{F}_n] = E[Z_{m,n+1}^{N_{max}}|X'_n]$. Next, note that by (1), $P[X'_j = 0|X'_n = 0] = 1$ for all $j = n + 1, n + 2, \dots, N_{max}$. Then by definition, $Z_{1,N_{max}}^{N_{max}} = 0$ and $P[Z_{m,r}^{N_{max}} = 0|X'_n = 0] = 1$ for any $1 \leq m \leq N_{max} - r + 1$, $r = n + 1, n + 2, \dots, N_{max} - 1$, and hence, $E[Z_{m,n+1}^{N_{max}}|X'_n = 0] = 0$. Also note that $E[Z_{m,n+1}^{N_{max}}|X'_{n+1} = 0] = 0$. Therefore, conditioning on X_{n+1} , $E[Z_{m,n+1}^{N_{max}}|X'_n > 0]$ can be expressed as

$$\begin{aligned}
E[Z_{m,n+1}^{N_{max}}|X'_n > 0] &= E[Z_{m,n+1}^{N_{max}}|X'_{n+1} > 0]P(X'_{n+1} > 0|X'_n > 0) \\
&+ E[Z_{m,n+1}^{N_{max}}|X'_{n+1} = 0]P(X'_{n+1} = 0|X'_n > 0) = \frac{\sum_{i=n+1}^{N_{max}} P_i}{\sum_{i=n}^{N_{max}} P_i} E[Z_{m,n+1}^{N_{max}}|X'_{n+1} > 0].
\end{aligned} \tag{EC.2}$$

By definition, $Z_{m,n+1}^{N_{max}}$ can be expressed through $Z_{m,n+2}^{N_{max}}$ and $Z_{m+1,n+2}^{N_{max}}$, and hence, (EC.2) becomes

$$\begin{aligned}
&\frac{\sum_{i=n+1}^{N_{max}} P_i}{\sum_{i=n}^{N_{max}} P_i} E[[X'_{n+1}|X'_{n+1} > 0] \vee E[Z_{m,n+2}^{N_{max}}|X'_{n+1} > 0] \wedge E[Z_{m-1,n+1}^{N_{max}}|X'_{n+1} > 0]] \\
&= \frac{\sum_{i=n+1}^{N_{max}} P_i}{\sum_{i=n}^{N_{max}} P_i} E[X_{n+1} \vee E[Z_{m,n+2}^{N_{max}}|X'_{n+1} > 0] \wedge E[Z_{m-1,n+1}^{N_{max}}|X'_{n+1} > 0]].
\end{aligned} \tag{EC.3}$$

Note that in (EC.3), $P[X'_{n+1} = X_{n+1}|X'_{n+1} > 0] = 1$ by (1). Since the distribution of X_{n+1} is given, then (EC.3) becomes (EC.1). \square

EC.3. Proof of Theorem 6.

THEOREM 6. *The optimal expected accumulated reward EV_1^C can be computed using the recursion,*

$$\begin{aligned}
 EV_j^c &= \frac{\sum_{i=j}^{N_{max}} P_i}{\sum_{i=j-1}^{N_{max}} P_i} [P(W_j \leq c, R_j + EV_{j+1}^{c-W_j} \geq EV_{j+1}^c) \\
 &\quad \times E[R_j + EV_{j+1}^{c-W_j} | W_j \leq c, R_j + EV_{j+1}^{c-W_j} \geq EV_{j+1}^c] \\
 &\quad + [P(R_j + EV_{j+1}^{c-W_j} < EV_{j+1}^c, W_j \leq c) + P(W_j > c)] EV_{j+1}^c], \tag{EC.4}
 \end{aligned}$$

with boundary conditions $EV_j^c = 0$ for any c and $j \geq N_{max}$.

Proof: By the definition of EV_j^c and conditioning on the event that job j does arrive,

$$\begin{aligned}
 EV_j^c &\equiv E(V_j^c | \text{job } j-1 \text{ has arrived}) = E(V_j^c | \text{job } j \text{ arrives}) P(\text{job } j \text{ arrives} | \text{job } j-1 \text{ has arrived}) \\
 &\quad + E(V_j^c | \text{job } j \text{ does not arrive}) P(\text{job } j \text{ does not arrive} | \text{job } j-1 \text{ has arrived}).
 \end{aligned}$$

Note that $P(\text{job } j \text{ arrives} | \text{job } j-1 \text{ has arrived}) = \frac{\sum_{i=j}^{N_{max}} P_i}{\sum_{i=j-1}^{N_{max}} P_i}$. Also, if job j does not arrive, then no reward can be collected, and hence, $E(V_j^c | \text{job } j \text{ does not arrive}) = 0$. Then, (EC.4) follows from Theorem 5, since the optimal assignment policy, described by (3), is used for making each allocation decision. Since no job can arrive after N_{max} jobs, no reward can be collected, which establishes the boundary conditions. \square

EC.4. Proof of Theorem 7.

THEOREM 7. *Assume that*

$$B \equiv E[\sup_j \left| \frac{X_j}{W_j} \right|] < +\infty, \text{ and } P[N < +\infty] = 1. \tag{EC.5}$$

Then for any $j = 1, 2, \dots$ and $c \in [0, C]$, the infinite sequence $\{EV_j^c(N_{max})\}_{N_{max}=1}^{+\infty}$ converges to the finite limit $EV_j^c(\infty) \equiv \lim_{N_{max} \rightarrow +\infty} EV_j^c(N_{max})$, and Theorem 5 establishes an optimal policy for DSKP, where the pmf of the number of jobs has infinite support, with EV_j^c replaced by $EV_j^c(\infty)$, $j = 1, 2, \dots$, $c \in [0, C]$.

Proof: For any $j = 1, 2, \dots$, $c \in [0, C]$, and $N_{max} = 1, 2, \dots$, $EV_j^c(N_{max} + 1) \geq EV_j^c(N_{max})$ and $EV_j^c(N_{max}) \leq Bc$. The infinite sequence $\{EV_j^c(N_{max})\}_{N_{max}=1}^{+\infty}$ is monotonically increasing and uniformly bounded, and hence, $\lim_{N_{max} \rightarrow +\infty} EV_j^c(N_{max})$ exists and is finite.

For any $j = 1, 2, \dots$ and $c \in [0, C]$, let OV_j^c denote the optimal *conditional* expected accumulated reward from the allocation of resource capacity c to jobs $j, j + 1, \dots$, given that job $j - 1$ has arrived. Theorem 5 establishes an optimal policy for DSKP, where the *pmf* of the number of jobs has infinite support, with EV_j^c replaced by OV_j^c , $j = 1, 2, \dots$, $c \in [0, C]$. It suffices to show that for any $j = 1, 2, \dots$ and $c \in [0, C]$, $OV_j^c = EV_j^c(\infty) \equiv \lim_{N_{max} \rightarrow +\infty} EV_j^c(N_{max})$.

For any $j = 1, 2, \dots$, $c \in [0, C]$ and $N_{max} = 1, 2, \dots$, $OV_j^c \geq EV_j^c(N_{max})$ and $\lim_{N_{max} \rightarrow +\infty} EV_j^c(N_{max}) \geq EV_j^c(\infty)$, and hence, $OV_j^c - EV_j^c(\infty) > 0$.

Suppose that an optimal policy for accepting jobs is followed up until the arrival of some job N , and then the jobs $N + 1, N + 2, \dots$ can be accepted without regard to their weights. Then, the expected total accumulated reward is higher than OV_j^c , and the following expression holds,

$$OV_j^c - EV_j^c(\infty) \leq EV_j^c(N) + \sum_{i=N+1}^{\infty} [X_i P(\text{job } i \text{ arrives})] - EV_j^c(\infty) \leq BC \sum_{i=N+1}^{\infty} [P(\text{job } i \text{ arrives})].$$

By (EC.5), for any $\Delta > 0$, there exists N_{Δ} such that $\sum_{i=N_{\Delta}+1}^{\infty} [P(\text{job } i \text{ arrives})] < \frac{\Delta}{BC}$, which means that with $N = N_{\Delta}$, $OV_j^c - EV_j^c(\infty) < \Delta$. Since Δ can be chosen arbitrarily, then $OV_j^c = EV_j^c(\infty)$.

□